

Ω -algebras

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Abstract

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Starting with Ω -sets where Ω is a complete lattice, we introduce the notion of an Ω -algebra. This is a classical algebra equipped with an Ω -valued equality replacing the ordinary one. In these new structures identities hold as appropriate lattice-theoretic formulas. Our investigation is related to weak congruences of the basic algebra to which a generalized equality is associated. Namely every Ω -algebra uniquely determines a closure system in the lattice of weak congruences of the basic algebra. By this correspondence we formulate a representation theorem for Ω -algebras.

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Ω -sets have been further applied to non-classical predicate logics, and also partially in theoretical foundations of fuzzy set theory.

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In our approach, an identity holds if the corresponding lattice-theoretic formula is fulfilled. An identity may hold on a lattice-valued algebra, while the underlying classical algebra need not satisfy the same identity.

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The collection $\text{Con}_w(\mathcal{A})$ of all weak congruences on an algebra \mathcal{A} is an algebraic lattice under inclusion.

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For $p \in L$, a **cut set** or a **p -cut** of an Ω -valued function $\mu : A \rightarrow \Omega$ is a subset μ_p of A which is the inverse image of the principal filter $\uparrow p$ in Ω : $\mu_p = \mu^{-1}(\uparrow p) = \{x \in X \mid \mu(x) \geq p\}$.

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An **Ω -valued (binary) relation** R on A is an Ω -valued function on A^2 , i.e., it is a mapping $R : A^2 \rightarrow \Omega$. As above, for $p \in \Omega$, a **cut** R_p of R is the binary relation on A , which is the inverse image of $\uparrow p$: $R_p = R^{-1}(\uparrow p)$.

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Strictness can be understood as a *weak reflexivity* of R . Therefore, a symmetric and transitive Ω -valued relation on A is a **weak Ω -valued equivalence on A** .

If $\mu : A \rightarrow \Omega$ is an Ω -valued function on A , then the map $R : A^2 \rightarrow \Omega$ on A is an **Ω -valued relation on μ** if for all $x, y \in A$

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*If $R : A^2 \rightarrow \Omega$ is a weak Ω -valued equivalence on A , then it is an Ω -valued equivalence on $\mu : A \rightarrow \Omega$, such that $\mu(x) = R(x, x)$. The Ω -valued function μ is said to be **determined** by R .*

A weak Ω -valued equivalence R on A is a weak **Ω -valued equality**, if it satisfies the **separation property**:

If $R(x, x) \neq 0$, then $R(x, y) = R(x, x)$ implies $x = y$.

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Analogously, an Ω -valued equivalence on $\mu : A \rightarrow \Omega$ satisfying the separation property is an **Ω -valued equality** on μ .

If $A = (A, F)$ is an algebra and $\mu : A \rightarrow \Omega$ an Ω -valued function on A , then μ is **compatible** with the operations in F , if for every n -ary operation $f \in F$, for all $a_1, \dots, a_n \in A$, and for every constant (nullary operation) $c \in F$

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Analogously, an Ω -valued relation $R : A^2 \rightarrow \Omega$ on A is **compatible** with the operations in F if for every n -ary operation $f \in F$, for all $a_1, \dots, a_n, b_1, \dots, b_n \in A$, and for every constant $c \in F$

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Similarly, an Ω -valued relation $R : A^2 \rightarrow \Omega$ on A is compatible with all the operations in F , if and only if for every $p \in \Omega$, R_p is compatible with all the operations in F .

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Lemma

Every cut $E_p = E^{-1}(\uparrow p)$, $p \in \Omega$, of the Ω -valued equality E in an Ω -set (A, E) is an equivalence relation on the corresponding cut μ_p of μ .

A pair $\overline{\mathcal{A}} = (\mathcal{A}, E)$ is an Ω -algebra if $\mathcal{A} = (A, F)$ is an algebra, (A, E) is an Ω -set and E is compatible with the operations in F . \mathcal{A} is the **underlying, basic algebra** of $\overline{\mathcal{A}}$.

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- (iii) Every cut of E is a weak congruence on \mathcal{A} , namely for $p \in E$, E_p is a congruence on μ_p .*

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Let (\mathcal{A}, E) be an Ω -algebra and $u(x_1, \dots, x_n) \approx v(x_1, \dots, x_n)$, briefly $u \approx v$ be an identity in the type of \mathcal{A} . We assume, as usual, that variables appearing in terms u and v are from x_1, \dots, x_n .

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Then, (\mathcal{A}, E) **satisfies identity** $u \approx v$ (i.e., this identity **holds** on (\mathcal{A}, E)) if

$$\bigwedge_{i=1}^n \mu(a_i) \leq E(u(a_1, \dots, a_n), v(a_1, \dots, a_n)),$$

for all $a_1, \dots, a_n \in A$ and the term-operations on \mathcal{A} corresponding to terms u and v respectively.

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Theorem

Let (\mathcal{A}, E) be an Ω -algebra, and \mathcal{F} a set of identities in the language of \mathcal{A} . Then, (\mathcal{A}, E) satisfies all identities in \mathcal{F} if and only if for every $p \in L$ the quotient algebra μ_p/E_p satisfies the same identities.

Corollary

If a diagonal relation $\Delta_A = \{(a, a) \mid a \in A\}$ is a cut of E , then each identity fulfilled by an Ω -algebra $\overline{\mathcal{A}} = (\mathcal{A}, E)$ also holds on the underlying algebra \mathcal{A} .

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By Corollary , we are interested in Ω -algebras which do not contain a copy of the underlying algebra among quotient substructures. An Ω -algebra $\overline{\mathcal{A}} = (\mathcal{A}, E)$ is said to be **proper** if Δ_A is not a cut of E .

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Theorem

$\bar{\mathcal{A}} = (\mathcal{A}, E)$ is a proper Ω -algebra if and only if

there are $a, b \in A$, $a \neq b$, such that $E(a, b) \geq \bigwedge \{E(x, x) \mid x \in A\}$.

Proposition

The collection of cuts of E in an Ω -algebra $\overline{\mathcal{A}} = (\mathcal{A}, E)$ is a closure system on A^2 , a subset of the weak congruence lattice $\text{Con}_w(\mathcal{A})$ of \mathcal{A} .

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Theorem (Representation)

Let \mathcal{A} be an algebra and \mathcal{R} a closure system in $\text{Con}_w(\mathcal{A})$ such that

if $a \neq b$, then $(a, b) \notin \bigcap \{R \in \mathcal{R} \mid (a, a) \in R\}$ for all $a, b \in A$.

Then there is a complete lattice Ω and an Ω -algebra (\mathcal{A}, E) with the underlying algebra \mathcal{A} , such that \mathcal{R} consists of cuts of E .

Sketch of the proof

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We take Ω to be the starting collection \mathcal{R} of weak congruences ordered by the dual of inclusion, \supseteq .

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Now we have that $E_R = R$ (the cut determined by R considered as an element of Ω , coincides with R as a weak congruence).

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Now we have that $E_R = R$ (the cut determined by R considered as an element of Ω , coincides with R as a weak congruence).

The structure (\mathcal{A}, E) is then the required Ω -algebra, obtained by the **canonical construction**.

For a symmetric and transitive relation $R \subseteq A^2$, we denote by $\text{dom}R$ the set $\{x \in A \mid (x, x) \in R\}$.

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Corollary

Let \mathcal{A} be an algebra and \mathcal{R} a closure system in $\text{Con}_w(\mathcal{A})$ fulfilling condition:

if $a \neq b$, then $(a, b) \notin \bigcap \{R \in \mathcal{R} \mid (a, a) \in R\}$ for all $a, b \in A$.

Let also \mathcal{F} be a set of identities in the language of \mathcal{A} and suppose that for every $R \in \mathcal{R}$, the algebra $\text{dom}R/R$ fulfills these identities. Then there is a complete lattice Ω and an Ω -algebra (\mathcal{A}, E) , such that \mathcal{R} consists of cuts of E and (\mathcal{A}, E) satisfies \mathcal{F} .

Suppose that we have different complete lattices, Ω_1 and Ω_2 and an algebra \mathcal{A} . Let $(\mathcal{A}, E1)$ and $(\mathcal{A}, E2)$ be an Ω_1 -valued algebra and an Ω_2 -valued algebra respectively. We say that the structures $(\mathcal{A}, E1)$ and $(\mathcal{A}, E2)$ are **cut-equivalent** if their collections of quotient algebras over cuts of $E1$ and $E2$ coincide, i.e., if for every $p \in \Omega_1$ there is $q \in \Omega_2$ such that $\mu_{1_p}/E1_p = \mu_{2_q}/E2_q$ and vice versa.

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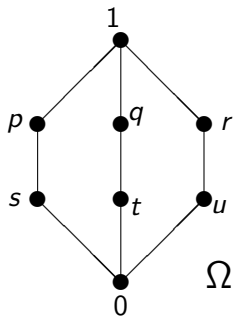
Theorem

Let $\overline{\mathcal{A}} = (\mathcal{A}, E)$ be an Ω -algebra where Ω is an arbitrary complete lattice. Then there is a lattice and a lattice-valued algebra cut-equivalent with $\overline{\mathcal{A}}$, obtained by the canonical construction over \mathcal{A} .

Examples

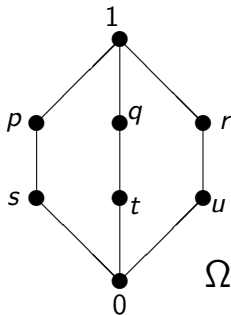
Examples

1. $((\{e, a, b, c, d\}, \cdot, '), E)$



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\cdot	e	a	b	c	d
e	e	a	b	c	d
a	a	e	d	a	a
b	b	c	c	e	a
c	c	c	e	b	b
d	d	d	d	d	e

	e	a	b	c	d
$'$	e	a	c	b	d

<i>E</i>	<i>e</i>	<i>a</i>	<i>b</i>	<i>c</i>	<i>d</i>
<i>e</i>	1	<i>u</i>	<i>t</i>	<i>t</i>	<i>s</i>
<i>a</i>	<i>u</i>	<i>r</i>	0	0	0
<i>b</i>	<i>t</i>	0	<i>q</i>	<i>t</i>	0
<i>c</i>	<i>t</i>	0	<i>t</i>	<i>q</i>	0
<i>d</i>	<i>s</i>	0	0	0	<i>p</i> .

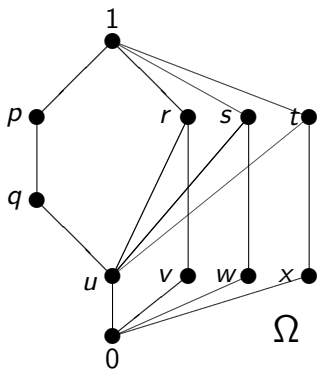
$$\mu = \begin{pmatrix} e & a & b & c & d \\ 1 & r & q & q & p \end{pmatrix}.$$

E	e	a	b	c	d
e	1	u	t	t	s
a	u	r	0	0	0
b	t	0	q	t	0
c	t	0	t	q	0
d	s	0	0	0	p .

$$\mu = \begin{pmatrix} e & a & b & c & d \\ 1 & r & q & q & p \end{pmatrix}.$$

The cuts of E are either diagonal relations on subalgebras (E_q on $\{e, b, c\}$ and E_r on $\{e, a\}$), or they are full relations on one-, two- or three-element subalgebras (e.g., E_t is a full relation on $\{e, b, c\}$). Trivially, E_0 is a full relation on the whole algebra. All the corresponding quotient algebras are groups, hence (\mathcal{A}, E) is an Ω -group. Observe that the basic five-element algebra is not a group.

2.



S_3 – symmetric group, (S_3, E) – the corresponding Ω -group.

\circ	e	f	g	h	j	k
e	e	f	g	h	j	k
f	f	e	h	g	k	j
g	g	j	e	k	f	h
h	h	k	f	j	e	g
j	j	g	k	e	h	f
k	k	h	j	f	g	e

E	e	f	g	h	j	k
e	1	x	w	q	q	v
f	x	t	u	0	0	u
g	w	u	s	0	0	u
h	q	0	0	p	q	0
j	q	0	0	q	p	0
k	v	u	u	0	0	r .

$$\mu = \begin{pmatrix} e & f & g & h & j & k \\ 1 & t & s & p & p & r \end{pmatrix}.$$

E	e	f	g	h	j	k
e	1	x	w	q	q	v
f	x	t	u	0	0	u
g	w	u	s	0	0	u
h	q	0	0	p	q	0
j	q	0	0	q	p	0
k	v	u	u	0	0	r .

$$\mu = \begin{pmatrix} e & f & g & h & j & k \\ 1 & t & s & p & p & r \end{pmatrix}.$$

All the structures μ_z/E_z , $z \in \Omega$ are groups of order 3, 2 or 1, hence Abelian.

E	e	f	g	h	j	k
e	1	x	w	q	q	v
f	x	t	u	0	0	u
g	w	u	s	0	0	u
h	q	0	0	p	q	0
j	q	0	0	q	p	0
k	v	u	u	0	0	r .

$$\mu = \begin{pmatrix} e & f & g & h & j & k \\ 1 & t & s & p & p & r \end{pmatrix}.$$

All the structures μ_z/E_z , $z \in \Omega$ are groups of order 3, 2 or 1, hence Abelian.

Therefore, this structure is an Abelian Ω -group, identity

$$x \cdot y \approx y \cdot x$$

holds as the formula

$$\mu(x) \wedge \mu(y) \leq E(x \cdot y, y \cdot x).$$

For the cuts, we have e.g., $\mu_p = \{e, h, j\}$, $\mu_u = \{e, f, g, h, j, k\}$.

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E_p	e	f	g	h	j	k
e	1	0	0	0	0	0
f	0	0	0	0	0	0
g	0	0	0	0	0	0
h	0	0	0	1	0	0
j	0	0	0	0	1	0
k	0	0	0	0	0	0

E_u	e	f	g	h	j	k
e	1	0	0	1	1	0
f	0	1	1	0	0	1
g	0	1	1	0	0	1
h	1	0	0	1	1	0
j	1	0	0	1	1	0
k	0	1	1	0	0	1

For the cuts, we have e.g., $\mu_p = \{e, h, j\}$, $\mu_u = \{e, f, g, h, j, k\}$.

E_p	e	f	g	h	j	k
e	1	0	0	0	0	0
f	0	0	0	0	0	0
g	0	0	0	0	0	0
h	0	0	0	1	0	0
j	0	0	0	0	1	0
k	0	0	0	0	0	0

E_u	e	f	g	h	j	k
e	1	0	0	1	1	0
f	0	1	1	0	0	1
g	0	1	1	0	0	1
h	1	0	0	1	1	0
j	1	0	0	1	1	0
k	0	1	1	0	0	1

Hence, E_p is a weak congruence on S_3 , a diagonal of $\mu_p = \{e, h, j\}$ and μ_p/E_p is a group of order 3.

For the cuts, we have e.g., $\mu_p = \{e, h, j\}$, $\mu_u = \{e, f, g, h, j, k\}$.

E_p	e	f	g	h	j	k
e	1	0	0	0	0	0
f	0	0	0	0	0	0
g	0	0	0	0	0	0
h	0	0	0	1	0	0
j	0	0	0	0	1	0
k	0	0	0	0	0	0

E_u	e	f	g	h	j	k
e	1	0	0	1	1	0
f	0	1	1	0	0	1
g	0	1	1	0	0	1
h	1	0	0	1	1	0
j	1	0	0	1	1	0
k	0	1	1	0	0	1

Hence, E_p is a weak congruence on S_3 , a diagonal of $\mu_p = \{e, h, j\}$ and μ_p/E_p is a group of order 3. Next, μ_u is the underlying group S_3 .

For the cuts, we have e.g., $\mu_p = \{e, h, j\}$, $\mu_u = \{e, f, g, h, j, k\}$.

E_p	e	f	g	h	j	k
e	1	0	0	0	0	0
f	0	0	0	0	0	0
g	0	0	0	0	0	0
h	0	0	0	1	0	0
j	0	0	0	0	1	0
k	0	0	0	0	0	0

E_u	e	f	g	h	j	k
e	1	0	0	1	1	0
f	0	1	1	0	0	1
g	0	1	1	0	0	1
h	1	0	0	1	1	0
j	1	0	0	1	1	0
k	0	1	1	0	0	1

Hence, E_p is a weak congruence on S_3 , a diagonal of $\mu_p = \{e, h, j\}$ and μ_p/E_p is a group of order 3. Next, μ_u is the underlying group S_3 . Therefore, $\mu_u/E_u = \{\{e, h, j\}, \{f, g, h\}\}$ i.e., it is a two-element quotient group, similarly for other cuts.

This Ω -group is obtained by the technique described above.

For the cuts, we have e.g., $\mu_p = \{e, h, j\}$, $\mu_u = \{e, f, g, h, j, k\}$.

E_p	e	f	g	h	j	k
e	1	0	0	0	0	0
f	0	0	0	0	0	0
g	0	0	0	0	0	0
h	0	0	0	1	0	0
j	0	0	0	0	1	0
k	0	0	0	0	0	0

E_u	e	f	g	h	j	k
e	1	0	0	1	1	0
f	0	1	1	0	0	1
g	0	1	1	0	0	1
h	1	0	0	1	1	0
j	1	0	0	1	1	0
k	0	1	1	0	0	1

Hence, E_p is a weak congruence on S_3 , a diagonal of $\mu_p = \{e, h, j\}$ and μ_p/E_p is a group of order 3. Next, μ_u is the underlying group S_3 . Therefore, $\mu_u/E_u = \{\{e, h, j\}, \{f, g, h\}\}$ i.e., it is a two-element quotient group, similarly for other cuts.

This Ω -group is obtained by the technique described above. The closure system i.e., the lattice Ω is $\text{Con}_w(S_3) \setminus \Delta_{S_3}$, consisting of all weak congruences on S_3 except the diagonal Δ_{S_3} . And the order in this lattice is dual to the set inclusion.

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Thanks for listening!